## Flag algebras in additive combinatorics

DOxML 2023 at GRIPS
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1. Additive Combinatorics ..... 4 slides
2. Constructive upper bounds through blow-ups ..... 3 slides3. Double counting on steroids5 slides
3. Conclusion ..... 1 slide

## Additive Combinatorics 101

Let $G$ be a finite (abelian) group of order $N$ or an interval $[N] \stackrel{\text { def }}{=}\{1, \ldots, N\} \subset \mathbb{N}$.

> What can we say about the (linear) structure of subsets $S \subseteq G$ or colorings $\gamma: G \rightarrow[c]$ ?

Global properties. What can we say about the relation of $|S|$ and Minkowski sumsets like $|S+S|$ ? Cauchy-Davenport, Vosper, Plünnecke-Ruzsa, Freiman-Ruzsa, ...

Local properties. Do $S$ or $G$ contain (monochromatic) $k$-term arithmetic progressions $(x, x+d, x+2 d, \ldots, x+(k-1) d)$, Schur triples $(x+y=z)$, repeated sums $(x+y=u+v)$ ? Schur, van der Waerden, Rado, Szémeredi, arithmetic regularity, Green-Tao, ...

## We will focus on the latter, in particular on the Rado Multiplicity Problem!

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## The Rado Multiplicity problem

Given a coloring $\gamma: G \rightarrow[c]$ and linear map $L: G^{m} \rightarrow G^{n}$, we are interested in

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\begin{equation*}
\mathcal{S}_{L}(\gamma, G) \stackrel{\text { def }}{=}\left\{\mathbf{s} \in G^{m}: L(\mathbf{s})=\mathbf{0}, s_{i} \neq s_{j} \text { for } i \neq j, \mathbf{s} \in \gamma^{-1}(\{i\})^{m} \text { for some } i\right\} \tag{1}
\end{equation*}
$$

Let $\Gamma_{c}(G)$ denotes all $c$-colorings of $G$. The Rado Multiplicity Problem is

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m_{q, c}(L, G) \stackrel{\text { def }}{=} \min _{\gamma \in \Gamma_{c}(G)}\left|\mathcal{S}_{L}(\gamma, G)\right| /\left|\mathcal{S}_{L}(G)\right|
$$

and in particular $m_{q, c}(L) \stackrel{\text { def }}{=} \lim \sup _{n \rightarrow \infty} m_{q, c}\left(L, G_{n}\right)$ when $G_{n}=[n], \mathbb{Z}_{n}, \mathbb{F}_{q}^{n}$.
Rado (1933) tells us that $\mathcal{S}_{L}\left(\gamma, G_{n}\right) \neq \emptyset$ if $L$ satisfies column condition and $n$ is large, which can also be shown to imply $0<m_{q, c}(L)<1$.

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## History of the problem

- Graham et al. (1996) gave lower bound for Schur triples in 2-colorings of [n], later independently resolved by Robertson and Zeilberger / Schoen / Datskovsky.
- Cameron et al. (2007) showed that the nr. of solutions for linear equations with an odd nr . of variables only depends on cardinalities of the two color classes.
- Parrilo, Robertson and Saracino (2008) established bounds for the minimum number of monochromatic 3-APs in 2-colorings of [n] (not 2-common in $\mathbb{N}$ ).
- For $r=1$ and $m$ even, Saad and Wolf (2017) showed that any 'pair-partitionable' $L$ gets its multiplicity from uniform random colorings of $\mathbb{F}_{q}^{n}$. Fox, Pham, and Zhao (2021) showed that this is necessary and Versteegen (2023) further generalized it.
- Kamčev et al. (2021) characterized some $L$ in $\mathbb{F}_{q}^{n}$ with $r>1$ where the multiplicity does not come from random constructions.
- Král et al. (2022) characterized $L$ where the mulitplicty comes from random constructions for $q=2, r=2, m$ odd.


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## Our results

We are interested in particular $L$ and $\mathbb{F}_{q}^{n}$.
Theorem (Rué and S., 2023)
We have $1 / 10<m_{q=5, c=2}\left(L_{4-\mathrm{AP}}\right) \leq 0.1 \overline{031746}$.
Saad and Wolf (2017) previously established an u.b. of 0.1247 with no no-trivial I.b. known.

Proposition (Rué and S., 2023)
We have $m_{q=3, c=3}\left(L_{3-A P}\right)=1 / 27$.
Similar to Cummings et al. (2013) extending a result of Goodman (1959) about triangles.

Both upper and lower bounds are computational in nature:

Upper bounds through blowup constructions of particular finite colorings. Discrete and Comb. Optimization

Lower bounds by extending Razborov's Flag Algebra framework. Conic Optimization, Sum-of-Squares, and Semidefinite Programming

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## 1. Additive Combinatorics

2. Constructive upper bounds through blow-ups
3. Double counting on steroids
4. Conclusion
5. Constructive upper bounds through blow-ups

## How to blow up colorings

What do we need for upper bounds? Sequences of colorings of increasing size... How can we turn this into a finite problem? By considering blowups:


Relevant in other contexts, e.g., Turán and Ramsey theory, capset problem, Sunflower conjecture, Turán's (3,4)-conjecture, the Shannon Capacity of odd cycles...

## Lemma

The limit of the density of monochromatic structures in the blow-up sequence is the non-injective density of monochromatic structures in the base coloring.
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## You can find constructions to blow up using your favorite Discrete Optimization technique:

isomorphism-free generation, SAT-solver, Integer Linear Programming, Bounded Tree Searches, Search Heuristics (Simmulated Annealing, Tabu Search, Genetic algorithms), even Machine Learning,

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## Proofs of the upper bounds

## Upper bound of the Theorem

 $m_{3,3}\left(L_{4-\mathrm{AP}}\right) \leq 1 / 27$ follows from the blow-up of this 3-coloring of $\mathbb{F}_{3}^{3}$ :

Upper bound of the Proposition $m_{5,2}\left(L_{4-\mathrm{AP}}\right) \leq 13 / 126$ follows from the iterated blow-up of this 2-coloring of $\mathbb{F}_{5}^{3}$ :


## 1. Additive Combinatorics

2. Constructive upper bounds through blow-ups

3 slides
3. Double counting on steroids 5 slides
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## An improvement on a trivial lower bound

The parameter $s_{L}(\gamma) \stackrel{\text { def }}{=}\left|\mathcal{S}_{L}(\gamma)\right| /\left|\mathcal{S}_{L}\left(\mathbb{F}_{q}^{n}\right)\right|$ satisfies the averaging equality

$$
\begin{equation*}
s_{L}(\gamma)=\sum_{\delta \in \Gamma(k)} p(\delta, \gamma) s_{L}(\delta)+o(1)=\mathbb{E}_{\delta \in \Gamma(k)}^{(\gamma)} s_{L}(\delta)+o(1) \tag{2}
\end{equation*}
$$

once $k$ is large enough. This implies an immediate trivial lower bound of

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\begin{equation*}
m_{q, c}(L) \geq \min _{\delta \in \Gamma(k)} s_{L}(\delta) \tag{3}
\end{equation*}
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If we magically found some coefficients $a_{\delta}$ satisfying $\mathbb{E}_{\delta \in \Gamma(k)}^{(\gamma)} a_{\delta}=o(1)$, we would get

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\begin{equation*}
m_{q, c}(L) \geq \min _{\delta \in \Gamma(k)} s_{L}(\delta)-a_{\delta} \tag{4}
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But how would we find such $a_{\delta}$ ? Flag Algebras and Semidefinite Programming!
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## 3. Double counting on steroids

## SOS please someone help me

## Definition

The flag algebra $\mathcal{A}$ is given by considering linear combinations of colorings, factoring out relations given by the averaging equality and defining an appropriate product.

The semantic cone $\mathcal{S}=\left\{f \in \mathcal{A}: \phi(f) \geq 0\right.$ for all $\left.\phi \in \operatorname{Hom}^{+}(\mathcal{A}, \mathbb{R})\right\}$ captures those algebraic expressions corresponding to density expressions that are 'true'.

There exists an element $C_{L} \in \mathcal{A}$ capturing the behavior of $S_{L}$, so we can establish a lower bound by establishing an SOS expression


The $p\left(\sum_{i=1}^{k} f_{i}^{2}, \delta\right)$ correspond to the $a_{\delta}$ on the previous slide!

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$$
\begin{equation*}
C_{L}-\lambda-\sum_{i=1}^{k} f_{i}^{2} \in \mathcal{S} \tag{5}
\end{equation*}
$$

The $p\left(\sum_{i=1}^{k} f_{i}^{2}, \delta\right)$ correspond to the $a_{\delta}$ on the previous slide!

## Challenges

1. Need an appropriate notion of density, isomorphism, and 'partially fixed coloring' both to (i) handle invariance and non-invariance and (ii) define different algebras.
2. Solutions as defined previously do not satisfy an exact averaging equality. $\rightarrow$ Introduce fully dimensional solutions, which asympt. make up all solutions.
3. Need to adequately solve isomorphisms problem from a practical perspective. $\rightarrow$ Represent structure as graph and use nauty.
4. (Almost) all SDP solvers work numerically, but we need algebraic expressions. $\rightarrow$ Refine solution an using exact LP solver like SoPlex.
5. Double counting on steroids

## Lower bound of the Proposition

$m_{5,2}\left(L_{4-\mathrm{AP}}\right)>1 / 10$ follows by verifying that over all 33242 -colorings of $\mathbb{F}_{5}^{2}$ we have

$$
\begin{aligned}
& F_{1}+F_{4}+\left(F_{2}+F_{3}\right) / 5-1 / 10 \geq \sum_{i=1}^{2}\left(9 / 10 \cdot \llbracket\left(F_{i, 1}+\left(5 F_{i, 2}-5 F_{i, 3}-10 F_{i, 4}\right) / 27\right)^{2} \rrbracket_{-1}\right. \\
&\left.\ldots+61 / 162 \cdot \llbracket\left(\left(F_{i, 3}-F_{i, 2}\right) / 2+F_{i, 4}\right)^{2} \rrbracket_{-1}\right),
\end{aligned}
$$

and by noting that $F_{1,1}+F_{2,1}>0$. Here the relevant flags $F_{i}$ and $F_{i, j}$ are

Flags of type $\varnothing$


Flags of type $\square$

$F_{1,4} \quad \square \square \square \square \square$

Flags of type

3. Double counting on steroids

## Lower bound of the Theorem

$m_{3,3}\left(L_{3-\mathrm{AP}}\right) \geq 1 / 27$ follows by verifying that over all all 140 3-colorings of $\mathbb{F}_{3}^{2}$ we have

$$
\begin{aligned}
F_{i}-1 / 27 \geq & 26 / 27 \cdot \llbracket\left(F_{i, 1}-99 / 182 F_{i, 2}+75 / 208 F_{i, 3}-11 / 28 F_{i, 4}-3 / 26 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+1685 / 1911 \cdot \llbracket\left(F_{i, 2}-231 / 26960 F_{i, 3}+1703 / 6740 F_{i, 4}-1869 / 3370 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+71779 / 431360 \cdot \llbracket\left(F_{i, 3}-358196 / 502453 F_{i, 4}-412904 / 502453 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+5431408 / 10551513 \cdot \llbracket\left(F_{i, 4}-1 / 4 F_{i, 5}\right)^{2} \rrbracket_{-1}
\end{aligned}
$$

for any $i \in\{1,2,3\}$. Here the relevant flags $F_{i}$ and $F_{i, j}$ are

| Flags of type $\varnothing$ | Flags of type $\square$ | Flags of type $\square$ | Flags of type $\square$ |
| :--- | :--- | :--- | :--- | :--- |
| $F_{1} \square \square \square$ | $F_{1,1} \square \square \square$ | $F_{2,1} \square \square \square$ | $F_{3,1} \square \square \square$ |
| $F_{2} \square \square \square$ | $F_{1,2} \square \square \square$ | $F_{2,2} \square \square \square$ | $F_{3,2} \square \square \square$ |
| $F_{3} \square \square \square$ | $F_{1,3} \square \square \square$ | $F_{2,3} \square \square \square$ | $F_{3,3} \square \square \square$ |
|  | $F_{1,4} \square \square \square$ | $F_{2,4} \square \square \square$ | $F_{3,4} \square \square \square$ |
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3 slides
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1 slide

## Open problems and final remarks

- Often one can extract stability results from Flag Algebra certificates.
- Steep computational hurdle: underlying structures grow exponentially (instead of quadratically for graphs or cubic for 3-uniform hypergraphs)
- No neat notion of subspaces makes generalizing to other groups difficult.

Code is available at github.com/FordUniver/rs_radomult_23

Thank you for your attention!

