

1. The Rado Multiplicity Problem
2. Lower bounds through Flag algebras in $\mathbb{F}_{q}^{n}$
3. Constructive upper bounds through blow-ups
4. Outlook

## Definition of the problem

Given a coloring $\gamma: \mathbb{F}_{q}^{n} \rightarrow[c]$ and linear map $L$, we are interested in

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\begin{equation*}
\mathcal{S}_{L}(\gamma) \stackrel{\text { def }}{=}\left\{\mathbf{s} \in\left(\mathbb{F}_{q}^{n}\right)^{m}: L(\mathbf{s})=\mathbf{0}, s_{i} \neq s_{j} \text { for } i \neq j, \mathbf{s} \in \gamma^{-1}(\{i\})^{m} \text { for some } i\right\} . \tag{1}
\end{equation*}
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Rado (1933) tells us that $\mathcal{S}_{L}(\gamma) \neq \emptyset$ for large enough $n$ if $L$ satisfies column condition.
The Rado Multiplicity Problem is concerned with determining

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m_{q, c}(L) \stackrel{\text { def }}{=} \lim _{n \rightarrow \infty} \min _{\gamma \in \Gamma(n)}\left|\mathcal{S}_{L}(\gamma)\right| /\left|\mathcal{S}_{L}\left(\mathbb{F}_{q}^{n}\right)\right| .
$$

Limit exists by monotonicity and $0<m_{q, c}(L) \leq 1$ if $L$ is partition regular. $L$ is $c$-common if $m_{q, c}(L)=c^{1-m}$ (the value attained in a uniform random coloring)

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## Previous results

- Graham et al. (1996) gave lower bound for Schur triples in 2-colorings of [ $n$ ], later independently resolved by Robertson and Zeilberger / Schoen / Datskovsky.
- Cameron et al. (2007) showed that the nr. of solutions for linear equations with an odd nr. of variables only depends on cardinalities of the two color classes.
- Parrilo, Robertson and Saracino (2008) established bounds for the minimum number of monochromatic 3-APs in 2-colorings of [n] (not 2-common in $\mathbb{N}$ ).
- For $r=1$ and $m$ even, Saad and Wolf (2017) showed that any 'pair-partitionable' $L$ is 2-common in $\mathbb{F}_{q}^{n}$. Fox, Pham, and Zhao (2021) showed that this is necessary.
- Kamčev et al. (2021) characterized some non-common $L$ in $\mathbb{F}_{q}^{n}$ with $r>1$.
- Král et al. (2022) characterized 2-common $L$ for $q=2, r=2$, $m$ odd.


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## What can we contribute?

We are interested in particular $L$ and fixed $q$. Analogous to e.g. determining Ramsey Multiplicity of $K_{4}$ for 2-uniform graphs (has a long history starting with Erdős in 1962).

## Theorem (Rué and S., 2023)

We have $1 / 10<m_{q=5, c=2}\left(L_{4-A P}\right) \leq 0.1031746$.

Saad and Wolf (2017) previously established an u.b. of 0.1247 with no no-trivial I.b. known.

Proposition (Rué and S., 2023)
We have $m_{q=3, c=3}\left(L_{3-A P}\right)=1 / 27$

Similar to Cummings et al. (2013) extending a result of Goodman (1959) about triangles.

Proofs are computational and applicable to other questions:

- Upper bounds through (iterated) blow-up constructions of particular finite colorings.
- Lower bounds through an extension of the Flag Algebra framework of Razborov (2007) to $\mathbb{F}_{q}^{n}$.


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## An improvement on a trivial lower bound

The parameter $s_{L}(\gamma) \stackrel{\text { def }}{=}\left|\mathcal{S}_{L}(\gamma)\right| /\left|\mathcal{S}_{L}\left(\mathbb{F}_{q}^{n}\right)\right|$ satisfies the averaging equality

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\begin{equation*}
s_{L}(\gamma)=\sum_{\delta \in \Gamma(k)} p(\delta, \gamma) s_{L}(\delta)+o(1)=\mathbb{E}_{\delta \in \Gamma(k)}^{(\gamma)} s_{L}(\delta)+o(1) \tag{2}
\end{equation*}
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once $k$ is large enough. This implies an immediate trivial lower bound of

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\begin{equation*}
m_{q, c}(L) \geq \min _{\delta \in \Gamma(k)} s_{L}(\delta) \tag{3}
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If we magically found some coefficients $a_{\delta}$ satisfying $\mathbb{E}_{\delta \in \Gamma(k)}^{(\gamma)} a_{\delta}=o(1)$, we would get

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m_{q, c}(L) \geq \min _{\delta \in \Gamma(k)} s_{L}(\delta)-a_{\delta} \tag{4}
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But how would we find such $a_{\delta}$ ? Flag Algebras and Semidefinite Programming!
2. Lower bounds through Flag algebras in $\mathbb{F}_{q}^{n}$

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Definition
The flag algebra (of the empty type) $\mathcal{A}$ is given by considering $\mathbb{R} \Gamma$, factoring out relations given by the averaging equality and defining an appropriate product.

There exists an element $C_{L} \in \mathcal{A}$ capturing the behavior of $s_{L}$ and the semantic cone

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\begin{equation*}
\mathcal{S}=\left\{f \in \mathcal{A}: \phi(f) \geq 0 \text { for all } \phi \in \operatorname{Hom}^{+}(\mathcal{A}, \mathbb{R})\right\} \tag{5}
\end{equation*}
$$

captures those algebraic expressions corresponding to density expressions that are 'true'. We can therefore establish a lower bound for $s_{L}$ by showing that


Here $p\left(\sum_{i=1}^{k}\left(f_{i}\right)^{2}, \delta\right)$ corresponds to the $a_{\delta}$ on the previous slide! Such sum-of-squares (SOS) expressions are related to easily solvable Semidefinite Programs (SDPs).

## The SDP approach

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\begin{equation*}
C_{L}-\lambda-\sum_{i=1}^{k}\left(f_{i}\right)^{2} \in \mathcal{S} \quad \Rightarrow \quad C_{L}-\lambda \in \mathcal{S} \tag{6}
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## Challenges

- Need an appropriate notion of density, isomorphism, and 'partially fixed coloring' both to (i) handle invariance and non-invariance and (ii) define different algebras.
- Solutions as defined previously do not satisfy an exact averaging equality. Need to introduce fully dimensional solutions, which asymptotically make up all solutions.
- Need to adequately solve isomorphisms problem from a practical perspective.
- (Almost) all SDP solvers work numerically, but we need algebraic expressions.

2. Lower bounds through Flag algebras in $\mathbb{F}_{q}^{n}$

## Lower bound of the Proposition

$m_{5,2}\left(L_{4-\mathrm{AP}}\right)>1 / 10$ follows by verifying that over all 3324 2-colorings of $\mathbb{F}_{5}^{2}$ we have

$$
\left.\begin{array}{rl}
F_{1}+F_{4}+\left(F_{2}+F_{3}\right) / 5-1 / 10 \geq & \sum_{i=1}^{2}\left(9 / 10 \cdot \llbracket\left(F_{i, 1}+\left(5 F_{i, 2}-5 F_{i, 3}-10 F_{i, 4}\right) / 27\right)^{2} \rrbracket_{-1}\right. \\
& \ldots+61 / 162 \cdot \llbracket\left(\left(F_{i, 3}-F_{i, 2}\right) / 2+F_{i, 4}\right)^{2} \rrbracket \\
-1
\end{array}\right),
$$

and by noting that $F_{1,1}+F_{2,1}>0$. Here the relevant flags $F_{i}$ and $F_{i, j}$ are

Flags of type $\varnothing$


Flags of type $\square$


Flags of type

2. Lower bounds through Flag algebras in $\mathbb{F}_{q}^{n}$

## Lower bound of the Theorem

$m_{3,3}\left(L_{3 \text {-AP }}\right) \geq 1 / 27$ follows by verifying that over all all 140 3-colorings of $\mathbb{F}_{3}^{2}$ we have

$$
\begin{aligned}
& F_{i}-1 / 27 \geq 26 / 27 \cdot \llbracket\left(F_{i, 1}-99 / 182 F_{i, 2}+75 / 208 F_{i, 3}-11 / 28 F_{i, 4}-3 / 26 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+1685 / 1911 \cdot \llbracket\left(F_{i, 2}-231 / 26960 F_{i, 3}+1703 / 6740 F_{i, 4}-1869 / 3370 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+71779 / 431360 \cdot \llbracket\left(F_{i, 3}-358196 / 502453 F_{i, 4}-412904 / 502453 F_{i, 5}\right)^{2} \rrbracket_{-1} \\
& \ldots+5431408 / 10551513 \cdot \llbracket\left(F_{i, 4}-1 / 4 F_{i, 5}\right)^{2} \rrbracket_{-1}
\end{aligned}
$$

for any $i \in\{1,2,3\}$. Here the relevant flags $F_{i}$ and $F_{i, j}$ are

| Flags of type $\varnothing$ | Flags of type $\square$ | Flags of type $\square$ | Flags of type $\square$ |
| :--- | :--- | :--- | :--- | :--- |
| $F_{1} \square \square \square$ | $F_{1,1} \square \square \square$ | $F_{2,1} \square \square \square$ | $F_{3,1} \square \square \square$ |
| $F_{2} \square \square \square$ | $F_{1,2} \square \square \square$ | $F_{2,2} \square \square \square$ | $F_{3,2} \square \square \square$ |
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## How to blow up colorings

We can blow up an colorings into a sequence of colorings with $n$ tending to infinity.


Computing the density of solutions in the limit of this sequence is easy: simply check not-necessarily-injective subcolorings in the base construction. This gives us an immediate upper bound from any coloring we can come up with ...

In some cases we have a free element in which we can iterate the blowup-construction.



## 3. Constructive upper bounds through blow-ups

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## 3. Constructive upper bounds through blow-ups <br> Proofs of the upper bounds

## Upper bound of the Proposition

$m_{5,2}\left(L_{4-\mathrm{AP}}\right) \leq 13 / 126$ follows from the iterated blow-up of this 2 -coloring of $\mathbb{F}_{5}^{3}$ :


## Upper bound of the Theorem

$m_{3,3}\left(L_{3 \text {-AP }}\right) \leq 1 / 27$ follows from the blow-up of this 3-coloring of $\mathbb{F}_{3}^{3}$

3. Constructive upper bounds through blow-ups

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## 4. Outlook

## Final Remarks

- Often one can extract stability results from Flag Algebra certificates.
- Steep computational hurdle: underlying structures grow exponentially (instead of quadratically for graphs or cubic for 3-uniform hypergraphs)

| $q / n$ | 1 | 2 | 3 | 4 | 5 |  | $q / n$ | 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 3 | 5 | 10 | 32 | 382 |  | 2 | 6 | 15 | 60 | 996 |
| 3 | 4 | 14 | 1028 |  |  |  | 3 | 10 | 140 | 25665178 |  |
| 4 | 8 | 1648 |  |  |  |  | 4 | 30 | 1630868 |  |  |
| 5 | 6 | 3324 |  |  |  |  | 5 | 24 | 70793574 |  |  |

Table: Number of 2- and 3-colorings of $\mathbb{F}_{q}^{n}$.

- No neat notion of subspaces makes generalizing to other groups difficult.

Code is available at github.com/FordUniver/rs_radomult_23

Thank you for your attention!

